

On the maximum size of arcs in the projective plane and some length optimal codes

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Projective plane $PG(2, q)$ over \mathbb{F}_q

Let \mathbb{F}_q be the finite field of order q . The set of n -tuples

$$\mathbb{F}_q^n = \{(x_1, x_2, \dots, x_n) \mid x_i \in \mathbb{F}_q\}$$

is the n -dimensional vector space over \mathbb{F}_q .

Let $PG(2, q)$ be the projective plane over \mathbb{F}_q , which consists of all lines through the origin of \mathbb{F}_q^3 over \mathbb{F}_q , that is,

$$(\mathbb{F}_q^3 \setminus \{(0, 0, 0)\}) / \sim,$$

where $(a, b, c) \sim (x, y, z) \iff (a, b, c) = \lambda(x, y, z)$ for some $\lambda \in \mathbb{F}_q \setminus \{0\}$.

$$PG(2, q) = \{(a, b, 1) \mid a, b \in \mathbb{F}_q\} \cup \{(a, 1, 0) \mid a \in \mathbb{F}_q\} \cup \{(1, 0, 0)\}.$$

The linear equation

is the equation of a line ℓ in $PG(2, q)$, simply denoted $\ell = [a, b, c]$, i.e.,

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(2) The set of points in $PG(2, 3)$ is

$$\{(0, 0, 1), (0, 1, 1), (0, 2, 1), (1, 0, 1), (1, 1, 1), (1, 2, 1), \\ (2, 0, 1), (2, 1, 1), (2, 2, 1), (0, 1, 0), (1, 1, 0), (2, 1, 0), (1, 0, 0)\},$$

and the set of lines in $PG(2, 3)$ is

$$\{[0, 0, 1], [0, 1, 1], [0, 2, 1], [1, 0, 1], [1, 1, 1], [1, 2, 1], \\ [2, 0, 1], [2, 1, 1], [2, 2, 1], [0, 1, 0], [1, 1, 0], [2, 1, 0], [1, 0, 0]\}.$$

The line $x + y + z = 0$ is

$$[1, 1, 1] := \{(0, 2, 1), (1, 1, 1), (2, 0, 1), (2, 1, 0)\}.$$

We note that $PG(2, 3)$ has 13 points and 13 lines and each line has 4 points in $PG(2, 3)$.

The following hold in $PG(2, q)$.

Points and lines in $PG(2, q)$

- ① $PG(2, q)$ consists of $q^2 + q + 1$ points and $q^2 + q + 1$ lines.
- ② Every line contains $q + 1$ points.
- ③ Two distinct lines meet at a point.
- ④ There are $q + 1$ lines passing through a point in $PG(2, q)$.

Consider quadrics in $PG(2, q)$

$$F := ax^2 + by^2 + cz^2 + dxy + eyz + fzx,$$

where $a, b, c, d, e, f \in \mathbb{F}_q$.

Define $v(F)$ as the zero set of F in $PG(2, q)$, i.e.,

$$v(F) := \{(x_0, x_1, x_2) \in PG(2, q) \mid F(x_0, x_1, x_2) = 0\}.$$

A conic means a nonsingular quadric.

Any conic has $q + 1$ (rational) points in $PG(2, q)$ with no three collinear.

A zero set of a singular quadric is a repeated line or a pair of distance lines or a point.

Arcs

An $(n, r)_q$ -arc

An $(n, r)_q$ -arc is a set \mathcal{K} of n points of $PG(2, q)$ such that some r but no $r + 1$ of them, are collinear, i.e., $|\mathcal{K} \cap \ell| \leq r$ for any line ℓ and $|\mathcal{K} \cap \ell| = r$ for some line ℓ in $PG(2, q)$.

For an $(n, r)_q$ -arc \mathcal{K} , the line ℓ is called i -line if $|\ell \cap \mathcal{K}| = i$. Define a_i as the number of i -lines to \mathcal{K} , i.e.,

$$a_i := \#\{\ell \mid |\ell \cap \mathcal{K}| = i\}.$$

Note that $a_i = 0$ for $i \geq r + 1$.

The $(r + 1)$ -tuple (a_0, a_1, \dots, a_r) is called the spectrum of the arc \mathcal{K} .

Example 2.

Let \mathcal{C} be a conic in $PG(2, q)$ with the equation $y^2 = xz$.

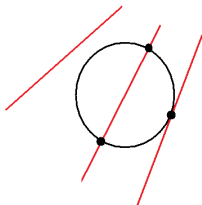
(1) The conic \mathcal{C} is a $(q+1, 2)_q$ -arc in $PG(2, q)$.

We note that the conic

$$\mathcal{C} = \{(t^2, t, 1) \mid t \in \mathbb{F}_q\} \cup \{(1, 0, 0)\}.$$

Thus $|\mathcal{C}| = q+1$ and for any line ℓ , we have $|\mathcal{C} \cap \ell| = 0$ or 1 or 2 .

Thus \mathcal{C} is a $(q+1, 2)_q$ -arc.



(2) For a conic \mathcal{C} , we have its spectrum as

$$a_0 = \frac{(q-1)q}{2}, \quad a_1 = q+1, \quad a_2 = \frac{(q+1)q}{2}.$$

Since the tangent lines of \mathcal{C} is 1-line, we have $a_1 = q+1$.

And each 2-line contains exactly two points of \mathcal{C} .

Thus we have

$$a_2 = \binom{q+1}{2} = \frac{(q+1)q}{2},$$

and hence

$$a_0 = q^2 + q + 1 - (a_1 + a_2) = \frac{(q-1)q}{2}.$$

Hence the spectrum is

$$a_0 = \frac{(q-1)q}{2}, \quad a_1 = q+1, \quad a_2 = \frac{(q+1)q}{2}.$$

The value of $m_r(2, q)$

Let $m_r(2, q)$ denote the largest n for which there exists an $(n, r)_q$ -arc for given r and q .

We call $(m_r(2, q), r)_q$ -arc the largest arc for given r and q .

An interesting problem in finite geometry is to determine the exact values of $m_r(2, q)$. Obviously, we have the bound for $m_r(2, q)$;

$$m_r(2, q) \leq (r - 1)q + r.$$

Linear codes

For two vectors $\mathbf{x} = (x_1, x_2, \dots, x_n)$ and $\mathbf{y} = (y_1, y_2, \dots, y_n)$ in \mathbb{F}_q^n , the Hamming distance between \mathbf{x} and \mathbf{y} , denoted by $d(\mathbf{x}, \mathbf{y})$ is the number of positions in which they differ.

An $[n, k, d]_q$ linear code C is a k -dimensional linear subspace of \mathbb{F}_q^n over \mathbb{F}_q with minimum distance d , where

$$\begin{aligned} d &= \min\{d(\mathbf{x}, \mathbf{y}) \mid \mathbf{x}, \mathbf{y} \in C, \mathbf{x} \neq \mathbf{y}\} \\ &= \min\{w(\mathbf{x}) \mid \mathbf{x} \in C, \mathbf{x} \neq \mathbf{0}\}. \end{aligned}$$

Here the weight $w(\mathbf{x})$ of \mathbf{x} is the number of nonzero positions in \mathbf{x} .

A good code will have small n (for fast transmission of messages), large k (for a wide variety of messages) and large d (to correct many errors).

Optimal linear codes problems

Optimize one of the parameters n , k , d for given values of the other two for a given field \mathbb{F}_q .

Optimal linear code problems

Optimal linear codes problems by Hill

- ① Find $B_q(n, d)$, the largest number q^k of codewords for which there exists an $[n, k, d]_q$ code.
- ② Find $d_q(n, k)$, the largest minimum distance d for which there exists an $[n, k, d]_q$ code.
- ③ Find $n_q(k, d)$, the smallest length n of codewords for which there exists an $[n, k, d]_q$ code.

A code which achieves one of the above values is called optimal, that is, we call dimension optimal, distance optimal and length optimal, respectively.

We note that the following;

A code is length optimal \implies distance optimal,
 \implies dimension optimal.

Example 1.

Let C_1 be a $[7, 2, 4]_2$ linear code with a generator matrix

$$\begin{pmatrix} 1 & 0 & 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}_{2 \times 7}.$$

Then C_1 is a distance optimal because there is no $[7, 2, 5]_2$ code. But C_1 is not a length optimal and a dimension optimal because there is a $[6, 2, 4]_2$ code and a $[7, 3, 4]_2$ code, respectively.

Griesmer bound

Griesmer bound

For an $[n, k, d]_q$ linear code, we have

$$n \geq g_q(k, d) := d + \left\lceil \frac{d}{q} \right\rceil + \left\lceil \frac{d}{q^2} \right\rceil + \cdots + \left\lceil \frac{d}{q^{k-1}} \right\rceil.$$

A linear code for which equality holds is called a Griesmer code.

We note that the Griesmer bound is an important lower bound of $n_q(k, d)$, that is,

$$n_q(k, d) \geq g_q(k, d).$$

Every Griesmer code is a length optimal.

Linear codes and arcs

Let C be an $[n, 3, d]_q$ code with a generator matrix G .
 Consider multi-set S whose elements are columns of G .
 Then S can be regarded as a multi-set in $PG(2, q)$.

Thus we have the following;

Theorem

An $[n, 3, d]_q$ linear code gives an $(n, n - d)_q$ -arc in $PG(2, q)$.

Equivalently, an $(n, r)_q$ -arc gives an $[n, 3, n - r]_q$ linear code.

Recall $m_r(2, q)$ denotes the largest arc for given r and q , and $m_r(2, q) \leq (r-1)q + r$.

Theorem.

For $(r-2)q + r < m_r(2, q) \leq (r-1)q + r$, the largest $(m_r(2, q), r)_q$ -arc corresponds to a Griesmer code (length optimal code).

We can easily see the following;

- ① For $r = 1$, the value $m_1(2, q) = 1$ and the arc is a singleton set.
- ② For $r = q$, the value $m_q(2, q) = q^2$ and the arc is the complement of a line ℓ_0 , i.e., $PG(2, q) \setminus \ell_0$.
- ③ For $r = q + 1$, the value $m_{q+1}(2, q) = q^2 + q + 1$ and the arc is the projective plane $PG(2, q)$.

A few values of $m_r(2, q)$, ($2 \leq r \leq q - 1$) are known for general q .

Known results on $m_r(2, q)$

Theorem. Bose (1947): On the values of $m_2(2, q)$

We have

$$m_2(2, q) = \begin{cases} q + 1, & q \text{ odd,} \\ q + 2, & q \text{ even.} \end{cases}$$

Theorem. Barlotti (1965) and Ball(1996)

For q odd prime, we have

$$m_r(2, q) = (r - 1)q + 1 \quad \text{for} \quad r = \frac{q + 1}{2} \quad \text{or} \quad r = \frac{q + 3}{2}.$$

Theorem. Denniston (1969)

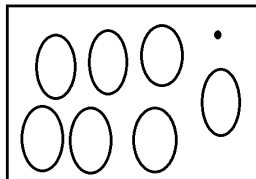
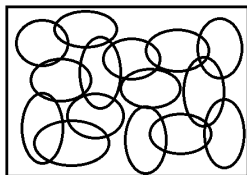
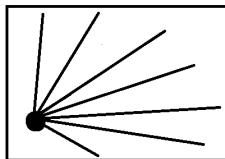
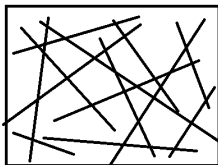
For q even, we have

$$m_r(2, q) = (r - 1)q + r \quad \text{for} \quad r = 2^e \leq q.$$

Problems

Let r be an integer with $2 \leq r \leq q - 1$.

- ① Find $m_r(2, q)$, the maximum value of n for which an (n, r) -arc exists in $PG(2, q)$.
Equivalently, find length optimal codes meeting the Griesmer codes.
- ② Classify (n, r) -arcs in $PG(2, q)$ for $n = m_r(2, q)$ up to projective equivalence.



A partition of $PG(2, q)$

For odd prime q , we denote by QR the set of quadratic residues mod q and by NQR the set of quadratic nonresidues mod q .

We consider a conic with the equation $ax^2 + by^2 + cz^2 + xy$ for $(a, b, c) \in PG(2, q)$, denoted by $f_{(a,b,c)}$.

Let $\mathcal{F} := \{f_{(a,b,c)} \mid (a, b, c) \in PG(2, q)\}$.

Lemma 1

For odd prime q , the conic $4xy - z^2$ and a line $[u, v, w]$ has no common point in $PG(2, q)$ if $w^2 - uv \in NQR$.

We express the line $[u, v, w]$ as follows;

$$[u, v, w] := \{(a, b, 1) \in PG(2, q) \mid au + bv + w = 0\} \cup \{(-v, u, 0)\}.$$

For two points $(a, b, 1), (a', b', 1) \in [u, v, w]$, consider two conics $C_1 := ax^2 + by^2 + z^2 + xy$ and $C_2 := a'x^2 + b'y^2 + z^2 + xy$. We have the following;

Lemma 2

Two conics C_1 and C_2 in \mathcal{F} which is given above are disjoint if $uv \in NQR$.

Next theorem shows that projective plane can be partitioned into disjoint conics and a point.

Theorem 3

For odd prime q , the projective plane $PG(2, q)$ consists of q disjoint conics and a point.

Proof) Choose q conics in \mathcal{F} such that $(a, b, c) \in [u, v, w]$ and $(a, b, c) \neq (-v, u, 0)$.

Here u, v and w satisfies that

$$uv \in NQR \text{ and } w^2 - uv \in NQR.$$

We note that such a line $[u, v, w]$ exists.

And the point $(0, 0, 1)$ is outside of the conics in \mathcal{F} . Thus we have

$$PG(2, q) = \{(0, 0, 1)\} \cup \bigcup_{(a,b,c) \in [u,v,w] \setminus \{(-v,u,0)\}} f_{(a,b,c)}.$$

The case $q = 7$

Consider conics in \mathcal{F} with the equation

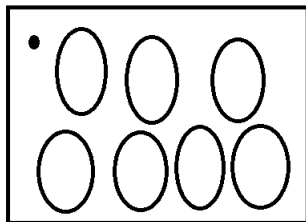
$$f_{(a,b,c)} = ax^2 + by^2 + cz^2 + xy,$$

and a line $[3, 1, 1]$. Then the line $[3, 1, 1]$ satisfies
 $uv = 3 \in NQR$ and $w^2 - uv = 1 - 3 \equiv -2 \equiv 5 \in NQR$.

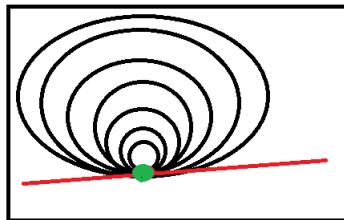
Choose 7 points (a, b, c)

on the line $[3, 1, 1]$, $(a, b, c) \neq (-1, 3, 0) = (2, 1, 0)$, where $[3, 1, 1] =$
 $\{(2, 0, 1), (0, 6, 1), (5, 5, 1), (1, 3, 1), (6, 2, 1), (4, 1, 1), (3, 4, 1), (2, 1, 0)\}$.

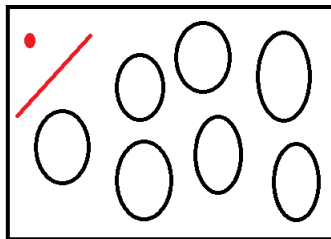
Then the union of 7 conics and a point $\{(0, 0, 1)\}$ is the whole
 plane $PG(2, 7)$.



Theorem 3



Theorem 4



Theorem 10

Conics with one common point

We consider the conics in \mathcal{F} with $f_{(a,b,c)}$, where (a, b, c) is on the line $[0, 1, 0] = \{(a, 0, 1) \mid a \in \mathbb{F}_q\} \cup \{(1, 0, 0)\}$.

Then the intersection of $f_{(a,b,c)}$, $(a, b, c) \in [0, 1, 0]$ is the point $(0, 1, 0)$.

Next theorem gives another geometrical configuration of $PG(2, q)$.

Theorem 4

For odd prime q , the projective plane $PG(2, q)$ consists of q conics with a common point P and common tangent line at P .

Proof) Consider the line $[0, 1, 0] = \{(a, 0, 1) \mid a \in \mathbb{F}_q\}$.

Choose q conics $f_{(a,0,1)}$ with $a \in \mathbb{F}_q, a \neq 1$.

Then the intersection of q conics $f_{(a,0,1)}, a \neq 1$ is the point $(0, 1, 0)$.

Note that the common tangent line of q conics $f_{(a,0,1)}$ is the line $[1, 0, 0]$. Then

$$PG(2, q) = [1, 0, 0] \cup \bigcup_{a \in \mathbb{F}_q, a \neq 1} f_{(a,0,1)}.$$



For $a \in \mathbb{F}_q \setminus \{0\}$, we have

$$\#\{a \in \mathbb{F}_q \mid 1 - a \in NQR, \quad a \neq 0\} = \frac{q-1}{2},$$

and

$$\#\{a \in \mathbb{F}_q \mid 1 - a \in NQR \cup \{0\}, \quad a \neq 0\} = \frac{q+1}{2}.$$

Using this, we have the following;

Theorem 5

(1) The union of $f_{(a,0,1)}$ satisfying $1-a \in NQR$ is $(\frac{q(q-1)}{2}+1, \frac{q+1}{2})_q$ maximal arc with the spectrum

$$a_0 = q, \quad a_1 = 1, \quad a_{\frac{q-1}{2}} = \frac{(q-1)q}{2}, \quad a_{\frac{q+1}{2}} = \frac{(q+1)q}{2}.$$

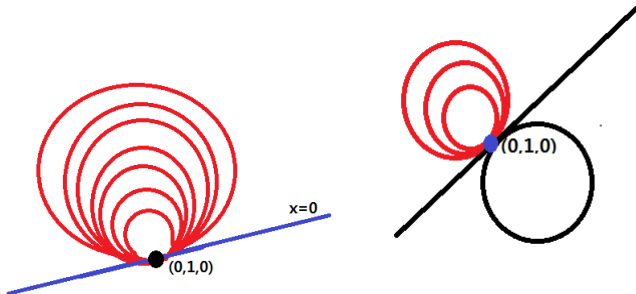
(2) The union of $f_{(a,0,1)}$ satisfying $1-a \in NQR \cup \{0\}$ is $(\frac{q(q+1)}{2}+1, \frac{q+3}{2})_q$ maximal arc with the spectrum

$$a_0 = 0, \quad a_1 = q+1, \quad a_{\frac{q+1}{2}} = \frac{(q-1)q}{2}, \quad a_{\frac{q+3}{2}} = \frac{(q+1)q}{2}.$$

Remark

The maximal arcs in Theorem 5 is exactly same arcs with one 's of Barlotti (1965) and Ball(1996). That is, for a given conic \mathcal{C} , if $t = \frac{q-1}{2}$, the arcs is the union of a point of \mathcal{C} and its internal set $\mathcal{I}(\mathcal{C})$ and $t = \frac{q+1}{2}$, the arcs is the union of the conic \mathcal{C} and its internal set $\mathcal{I}(\mathcal{C})$.

For $q = 7$, we have

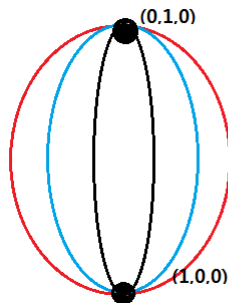


The case $q = 13$

In $PG(2, 13)$, consider three conics with two common points $(0, 1, 0)$ and $(1, 0, 0)$ as follows;

$$C_1 := 4z^2 + xy, \quad C_2 := 5z^2 + xy \quad C_3 := 12z^2 + xy.$$

Then the union of C_1, C_2 and C_3 is a $(38, 4)_{13}$ maximal arc.



The case $q = 8$

When q is even, we consider only $q = 8$.

It is known that the largest $(28, 4)_8$ -arc exists uniquely upto projective equivalence and it's spectrum is $a_0 = 10$ and $a_4 = 63$.

We represent that $(28, 4)_8$ -arc as several ways.

Let α be a primitive element of \mathbb{F}_8 with $\alpha^3 + \alpha + 1 = 0$.

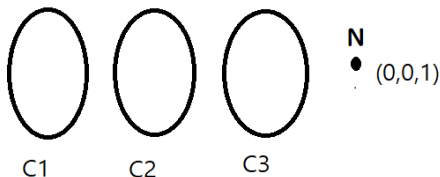
Here we express the arc with conics.

(1) The $(28, 4)_8$ -arc is expressed as the union of three disjoint conics with common nucleus.

Let C_i ($i = 1, 2, 3$) be a conic defined by the equation

$$C_i : x^2 + y^2 + \lambda_i z^2 + xy = 0,$$

where $\{\lambda_1, \lambda_2, \lambda_3\} = \{1, \alpha, \alpha^3\}$. Then the point $(0, 0, 1)$ is the nucleus of C_i ($i = 1, 2, 3$). The set $C_1 \cup C_2 \cup C_3 \cup \{(0, 0, 1)\}$ forms the $(28, 4)_8$ -arc.



(2) The $(28, 4)_8$ -arc is expressed using four conics passing through two points.

Let C_i ($i = 1, \dots, 4$) be a conic defined by the equation

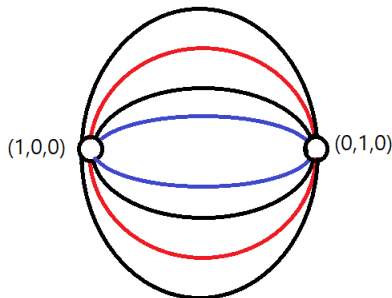
$$C_i : \mu_i z^2 + xy = 0,$$

where $\{\mu_1, \mu_2, \mu_3, \mu_4\} = \{1, \alpha^3, \alpha^5, \alpha^6\}$.

Then the point $(0, 0, 1)$ is the nucleus of C_i ($i = 1, \dots, 4$).

For $i \neq j$, $C_i \cap C_j = \{(1, 0, 0), (0, 1, 0)\}$.

The set $\cup_{i=1}^4 C_i \setminus \{(1, 0, 0), (0, 1, 0)\}$ is a $(28, 4)_8$ -arc.



(3) For a conic C in $PG(2, 8)$, the conic C has the spectrum
 $a_0 = 28$, $a_1 = 9$, $a_2 = 36$.

Then the number of 0-lines is 28.

Those lines forms a dual $(28, 4)_8$ -arc, i.e., the set
 $\{(a, b, c) \in PG(2, 8) \mid [a, b, c] \text{ is a 0-line of } C\}$ is the $(28, 4)_8$ -arc.

Zero set of a polynomial

Theorem 6. (Homma and Kim, 2018)

Let k be an integer with $1 \leq k \leq n$, and $PG(k-1, q)$ the linear subspace defined by $x_k = x_{k+1} = \cdots = x_n = 0$. Then the ideal of $PG(n, q) \setminus PG(k-1, q)$ in $\mathbb{F}_q[x_0, \dots, x_n]$ is generated by

$$\{x_i^q x_j - x_i x_j^q \mid 0 \leq i < j \leq n\} \\ \cup \left\{ x_s \prod_{i=k}^n (x_i^{q-1} - x_s^{q-1}) \mid s = 0, 1, \dots, k-1 \right\}.$$

Corollary 7.

The ideal of $PG(n, q) \setminus \{(1, 0, \dots, 0)\}$ in $\mathbb{F}_q[x_0, \dots, x_n]$ is generated by $\{x_i^q x_j - x_i x_j^q \mid 0 \leq i < j \leq n\} \cup \{x_0 \prod_{i=1}^n (x_i^{q-1} - x_0^{q-1})\}$.

Corollary 8.

Let $P_0 \in PG(n, q)$. Then there is a homogeneous polynomial F of degree d in $\mathbb{F}_q[x_0, \dots, x_n]$ such that the hypersurface H defined by $F = 0$ satisfies $H(\mathbb{F}_q) = PG(n, q) \setminus \{P_0\}$ if and only if $d \geq (q-1)n + 1$.

Corollary 9.

Let $P_0 \in PG(2, q)$. Then there is a homogeneous polynomial F of degree d in $\mathbb{F}_q[x_0, x_1, x_2]$ such that the curve H defined by $F = 0$ satisfies $H(\mathbb{F}_q) = PG(2, q) \setminus \{P_0\}$ if and only if $d \geq 2q - 1$.

Another partition of $PG(2, q)$ with conics

Next theorem gives another partition of $PG(2, q)$ with conics.

For odd q , the following holds.

Theorem 10.

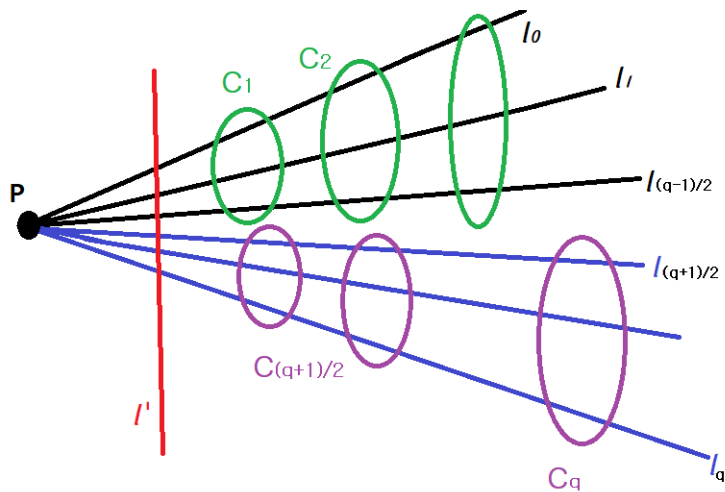
Let P be a point in $PG(2, q)$ and let ℓ_i ($i = 0, \dots, q$) be $q + 1$ lines passing through P . Then there exists disjoint $q - 1$ conics C_1, \dots, C_{q-1} , a line ℓ' and satisfying the following conditions;

$$PG(2, q) = \cup_{i=1}^{q-1} C_i \cup \ell' \cup \{P\}.$$

and

$$\cup_{i=1}^{\frac{q-1}{2}} C_i \subset \cup_{i=0}^{\frac{q-1}{2}} \ell_i \text{ and } \cup_{i=\frac{q+1}{2}}^{q-1} C_i \subset \cup_{i=\frac{q+1}{2}}^{q+1} \ell_i,$$

by renumbering the lines ℓ_i .



Let H be the set $(\cup_{i=2}^{\frac{q-1}{2}} C_i) \cup (\cup_{i=\frac{q+1}{2}}^{q+1} \ell_i) \cup \ell'$.

Then H is the complement of the conic C_1 , and of degree

$$(q-3) + \frac{q+1}{2} + 1 = \frac{3}{2}(q-1).$$

We obtained a polynomial H whose zero set is the complement of the conic C_1 .

Now we prove that the $\deg H$ is the minimum of such polynomials whose zero set is the complement of C_1 .

Theorem 11.

We have

$$\min\{\deg f \mid v(f) = PG(2, q) \setminus C_1\} = \frac{3}{2}(q - 1).$$





Proof. Suppose that there exists a polynomial f of degree $\leq \frac{3}{2}(q - 1) - 1$ such that $v(f) = PG(2, q) \setminus C_1$.

Let $\{Q_0, Q_1, \dots, Q_q\} = C_1$.

Consider $\frac{q-1}{2}$ lines $\overline{Q_1 Q_2}, \dots, \overline{Q_{q-2} Q_{q-1}}$ and any line ℓ containing Q_q but not Q_0 . Then the product of f and these $\frac{q+1}{2}$ lines is the polynomial of degree $\leq 2q - 2$, whose zero set is exactly $PG(2, q) \setminus \{Q_0\}$. It contradicts the Corollary 9. □

Thank you for your attention!!

References

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